

CMPT 478/981 Spring 2025

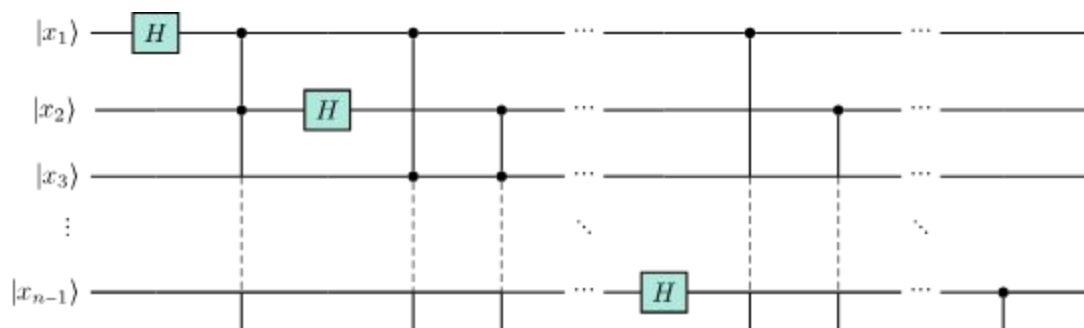
Quantum Circuits & Compilation

Matt Amy

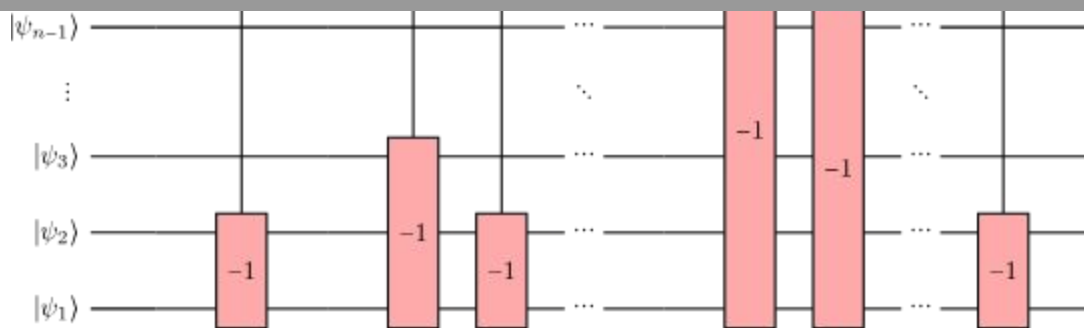
Today's agenda

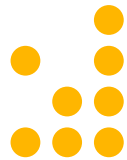


- Paper discussion (Classical oracles)
- Time evolution (“Quantum” oracles)



Time evolution





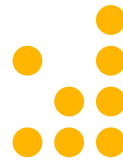
Hamiltonian simulation

- Quantum computing's “killer app”
- Idea is to simulate a quantum mechanical system with a quantum computer
- Works by implementing the **time evolution operator** $U(t) = e^{-iHt}$ where H is a Hermitian matrix called the system's **Hamiltonian**
- Compilation of $U(t)$ depends on the Hamiltonian but generic results exist

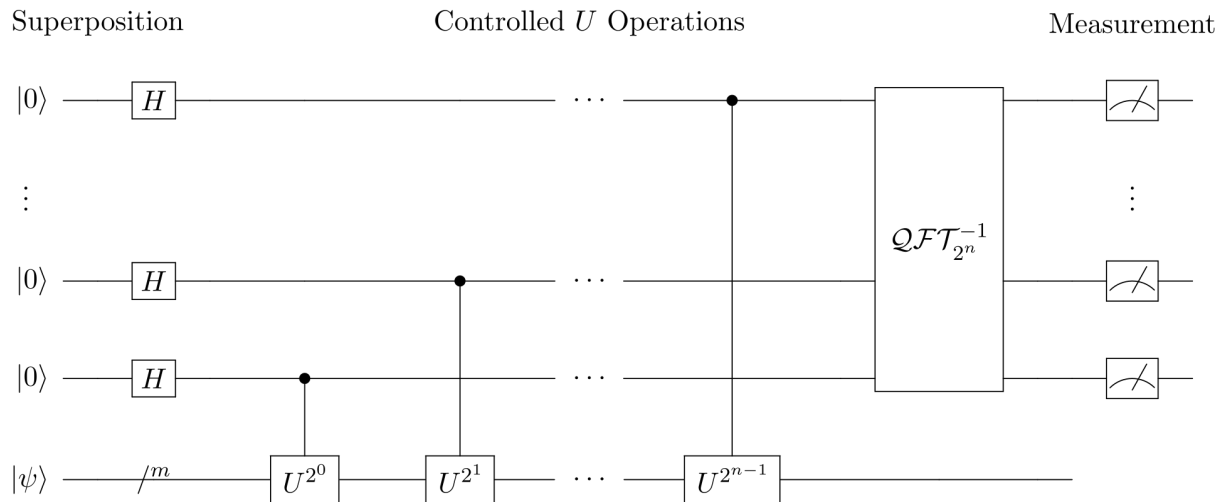
Theorem-ish (Lloyd, 1996)

A k -local Hamiltonian can be implemented in poly-time (ish)

Example: Ground state estimation



- Use phase estimation (i.e. Shor) to estimate an eigenvalue of Hamiltonian
- Replaces modular multiplication with time evolution operator $U(t)$



Implementing the time evolution operator

- In chemistry contexts, usually a **fermionic** Hamiltonian

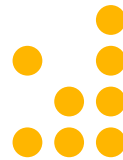
$$H = \sum_{p,q} h_{pq} \hat{a}_p^\dagger \hat{a}_q + \sum_{p,q,r,s} g_{pqrs} \hat{a}_p^\dagger \hat{a}_q^\dagger \hat{a}_r \hat{a}_s$$

- Apply a mapping to spin (Pauli) operators to get

$$H' = \sum_i s_i P_i + \sum_{i,j} d_{ij} P_i P_j + \dots$$

- **Problem: Pauli terms don't commute, i.e. $e^{A+B} \neq e^A e^B$**

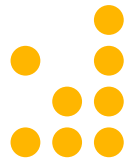
Classic solution: Product formulas



- Trotter formula gives limit of e^{A+B} as many alternations of fractions of $e^A e^B$

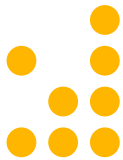
$$e^{i(A+B)t} = \lim_{n \rightarrow \infty} (e^{iA/n} e^{iB/n})^n$$

- Basic idea: select small enough n =epsilon to give good enough error bound
- Result is a **loooooong**, repeating string of Pauli exponentials around small angles



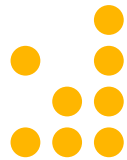
Pauli exponentials

- Recall: Cliffords map Paulis to Paulis
- Given e^{-iPt} for Pauli P , diagonalize as $Ce^{-i(I \otimes I \otimes \dots \otimes I \otimes Z)t}C^\dagger$
 - \Rightarrow Implement as Clifford basis change + single-qubit Z rotation



Pauli partitioning

- Important optimization in both NISQ & FTQEC regimes
- Problem: Given a spin Hamiltonian H , break up into minimal number of **commuting subsets** of Paulis
 - E.g. $H = ZXZ + IZI + IIX + XZX \rightarrow \{ZXZ\}, \{IZI, IIX, XZX\}$
- Commuting subset can be simultaneously diagonalized!
 - Map first Pauli to $ZI\dots I$
 - Remaining Paulis only have Z in first qubit, so iterate on next $n-1$ qubits
 - Result is a set of tensor products of Z
 - \rightarrow Exponential is a phase polynomial! More on optimizing those later...

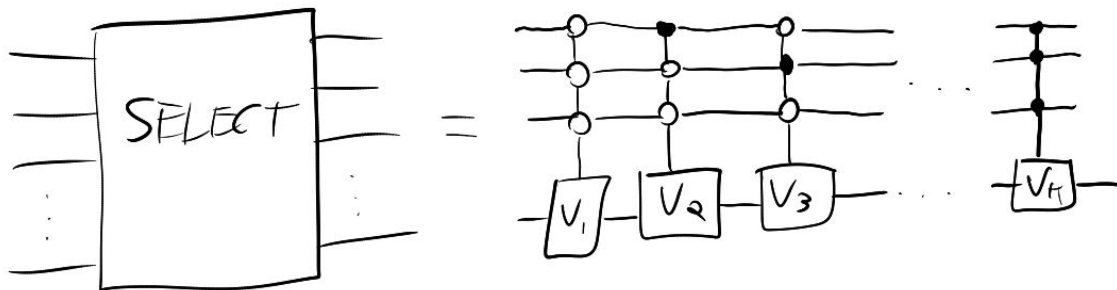


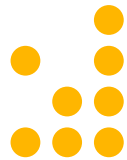
Implementing commuting Paulis

- Result is some set $\{P\}$ of Z-type Paulis
 - $e^{-iaZ_iZ_jZ_k}|x\rangle = e^{-ia(x_i \oplus x_j \oplus x_k)}|x\rangle$, i.e. Phase polynomial!
- How to iterate through linear functions of n bits?
 - Gray code!
- In general don't need all linear functions, so can we do better?
 - Heuristically, yes (Gray-synth, Amy Azimzadeh & Mosca "On the CNOT-complexity of CNOT-phase circuits")
 - NP-hard in special cases
 - Complexity of general case still unknown

Modern solutions: linear combinations of unitaries

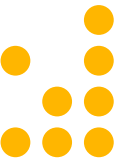
- Modern alternative to product formulas is to use Linear Combination of Unitaries (LCU) in various ways
- Basic idea: to implement the sum $aU + bV$ of unitaries
 - (PREPARE) Prepare a register in the state $a|0\rangle + b|1\rangle$
 - (SELECT) Apply a quantum multiplexor which sends $|0\rangle|\psi\rangle \rightarrow |0\rangle U|\psi\rangle$ and $|1\rangle|\psi\rangle \rightarrow |1\rangle V|\psi\rangle$





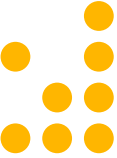
Hamiltonian simulation “algorithms”

- = methods of compiling a unitary circuit approximating time evolution
- Product-formula based
 - Product formulas
 - Multi-product formulas
 - Higher-order product formulas
 - QDrift
- LCU-based
 - Taylor series approximation
 - Quantum Signal Processing
 - Qubitization



Compiling time evolution operators

- Non-trivial task to automate due to
 - calculation of error bounds
 - calculation of e.g. roots of high degree polynomials (QSP)
 - combinatorial explosion of parameters & combinations of techniques
 - Jordan-Wigner vs Bravyi-Kitaev vs direct Fermionic Hamiltonian
 - ~6-7 algorithmic **frameworks** each with choices within
 - Possibility to combine techniques for separate parts of the Hamiltonian
 - Different models of Hamiltonians & qubit reduction techniques
- Goal for compilers is to automate at least some of the frequent tasks
 - LCU/block encodings
 - Grouping & compiling sequences of Pauli exponentials



Readings for next week

- Posted to the website
 - Childs et al., *Toward the first quantum simulation with quantum speedup*. arxiv.org:1711.10980
 - Focus on the main paper & understand the high-level structure of the various algorithms
 - You may skip over the derivation and analysis of error bounds
 - Nam, Su, Maslov, *Approximate Quantum Fourier Transform with $O(n\log(n))$ T gates*. arXiv:1803.04933
 - Pay attention here to the **use of an adder to implement the phase gates**
 - Campbell, *A random compiler for fast Hamiltonian simulation*. arXiv:1811.08017
- As before send me a short (paragraph or two) summary of **ONE (1)** paper of your choice before next class and be prepared to give a short summary of any of the papers in class